



# An epistemic logic for modeling decisions in the context of incomplete knowledge

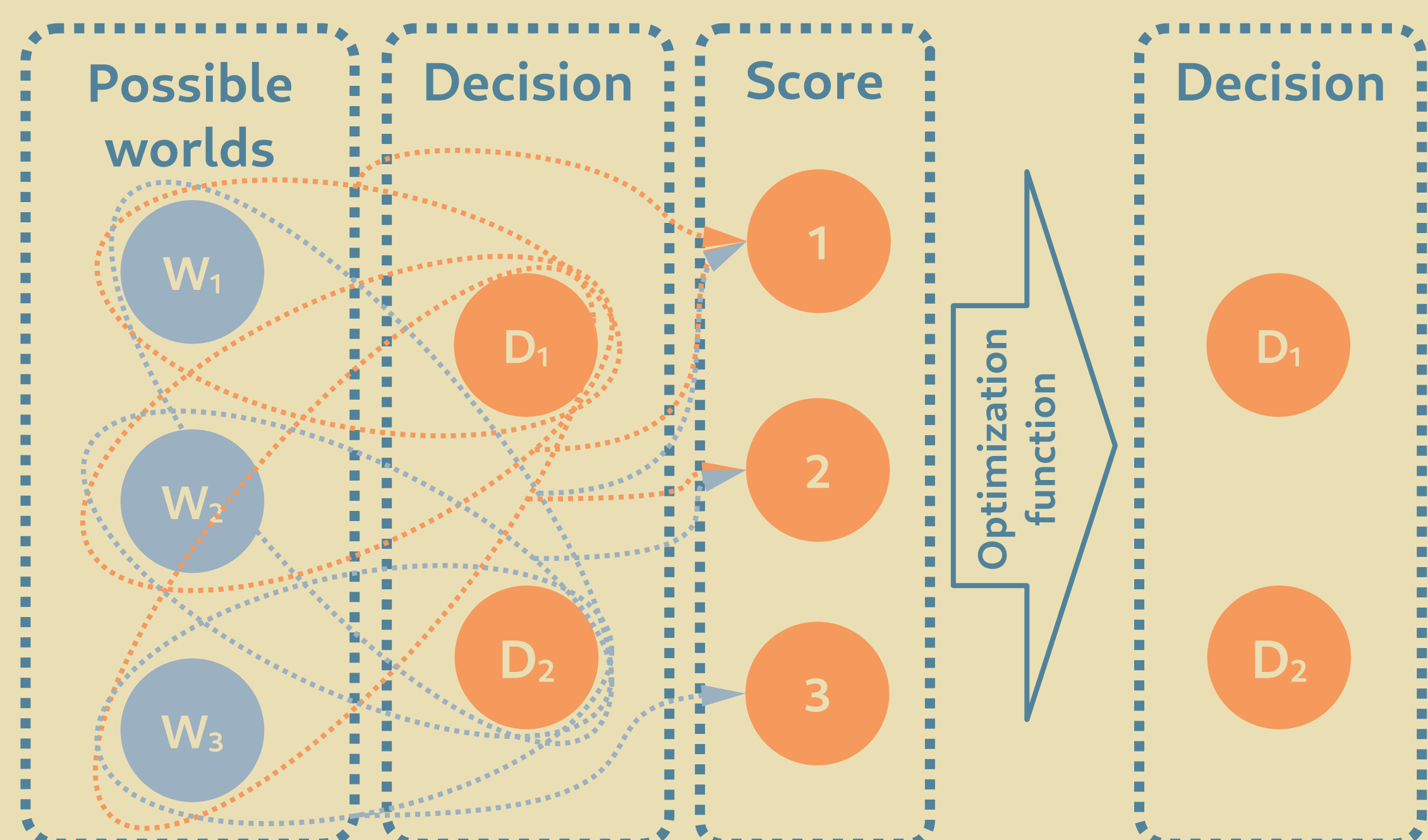
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## How do we model decision making processes?

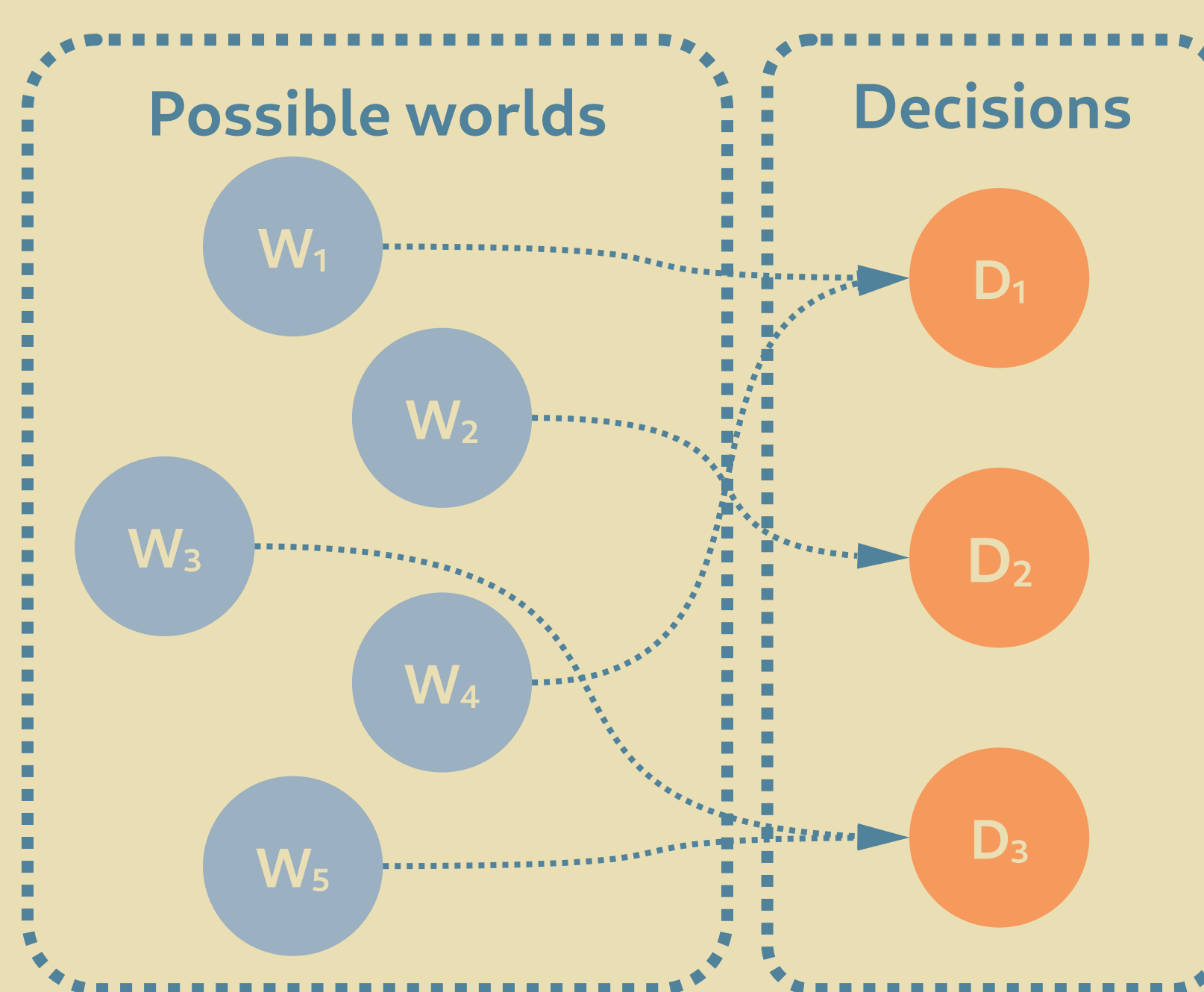
### Decision theory<sup>1</sup>



**Utility function** - mapping possible worlds and decisions to an ordinal. Then the decision is made based on the **preference relation**.

⚠ Too complex (modeling-wise)

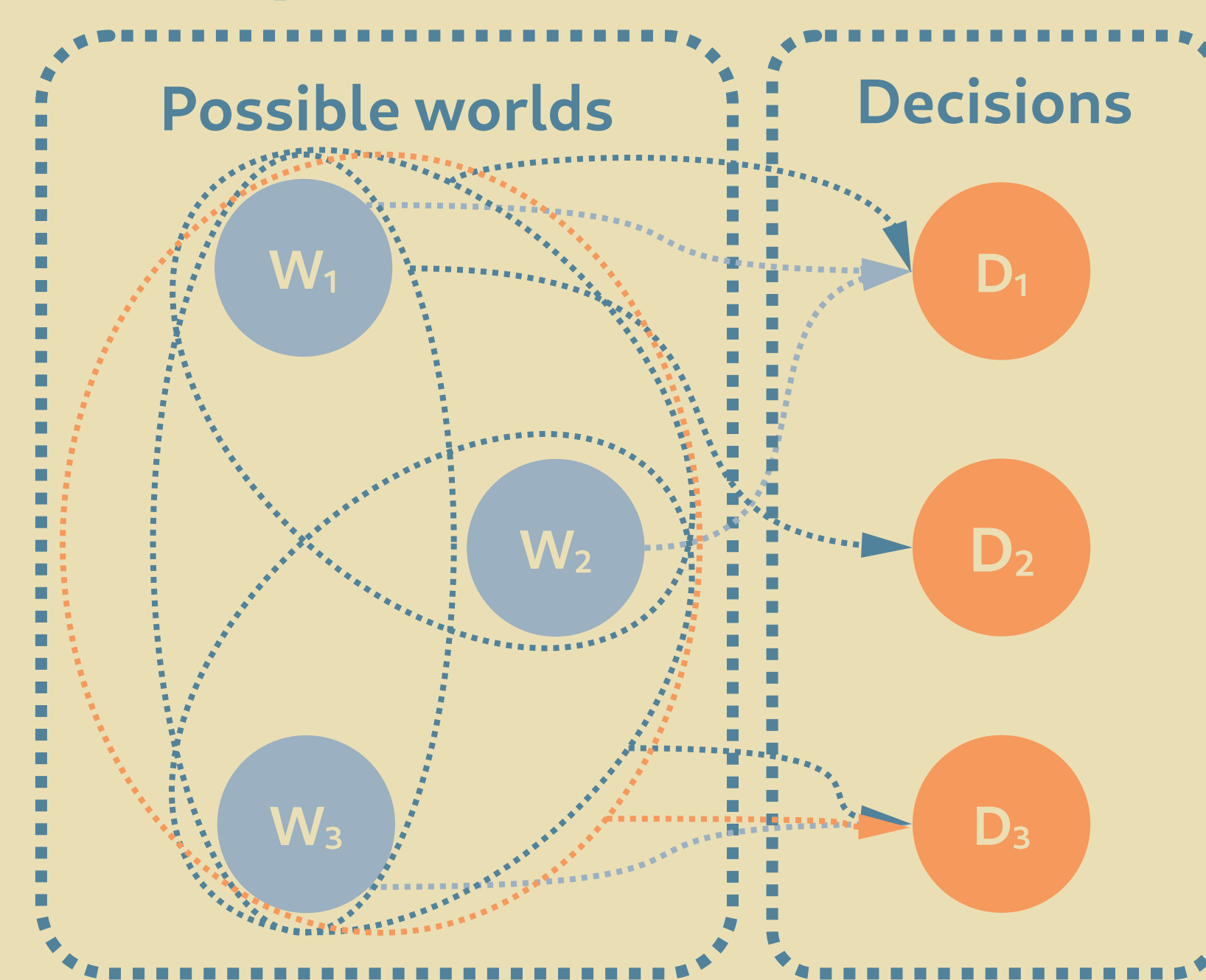
### Rule-based systems



Generalization of notations like **DMN** - objective states of affairs are directly related to decisions.

⚠ No support for incomplete knowledge!

### Epistemic decisions



**Epistemic decisions**<sup>1</sup> – mapping all possible epistemic states (sets of objective state of affairs) to decisions.

✓ Relation between knowledge and decision

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## Decision Model and Notation

### Customer greeting example:

| U | Gender | Marital status | Salutation |
|---|--------|----------------|------------|
| 1 | Male   | -              | Mr         |
| 2 | Female | Single         | Ms         |
| 3 | Female | Married        | Mrs        |

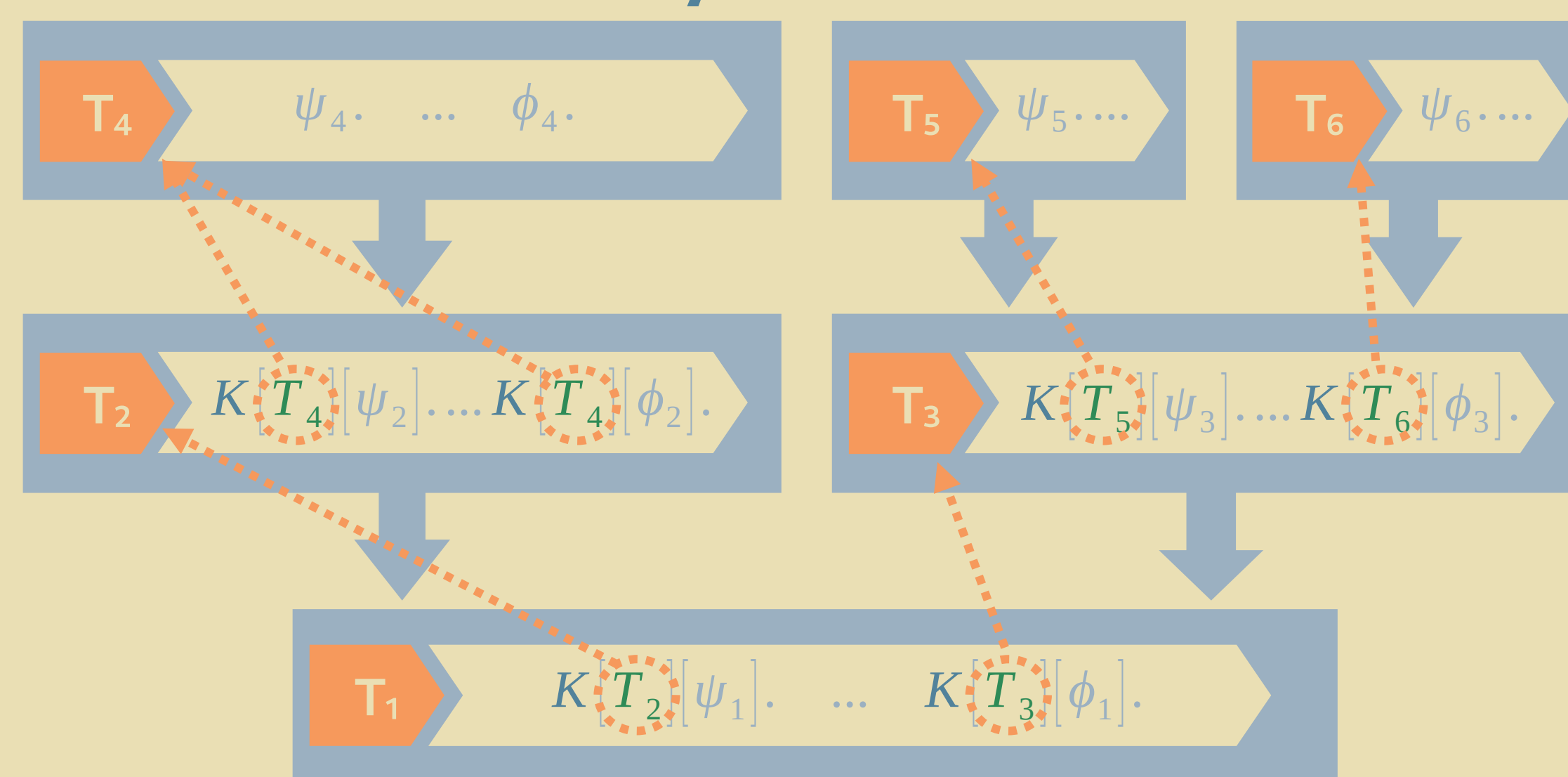
It is impossible to make a decision when some data is unknown.

- “-” stands for the absence of the constraint on the value.
- In DMN it is impossible to make a decision if some value is unknown: For example, the marital status.
- Practically, it is possible to make decision if gender is known to be “male” since it does not depend on marital status.

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## Ordered Epistemic Logic

### Hierarchy of theories



- $T_1 \dots T_6$  - FO(ID) Theories, over the same vocabulary.
- $K$  operator refers to other theories.
- $K$ -referencing is a partial order.
- $T_4 \dots T_6$  - Do not contain the  $K$  operator, they represent knowledge of objective state of affairs.

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## Result - epistemic DMN

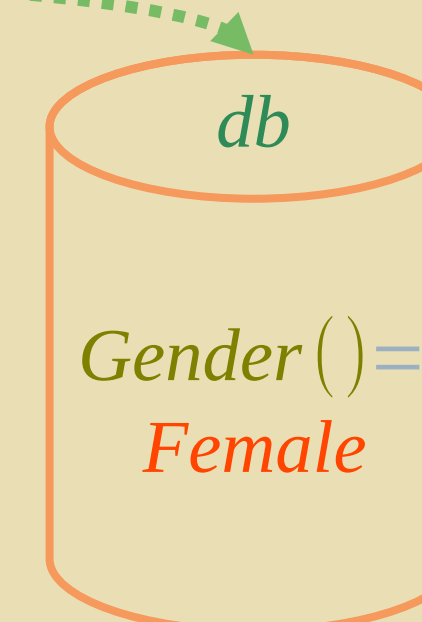
### Greeting example in eDMN

| U | Gender | Marital status | Salutation |
|---|--------|----------------|------------|
| 1 | Male   | -              | Mr         |
| 2 | Female | Single         | Ms         |
| 3 | Female | Married        | Mrs        |
| 4 | Female | ⊥              | Lady       |
| 5 | ⊥      | -              | Customer   |

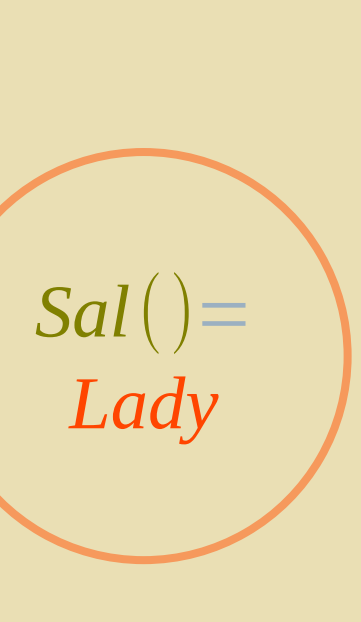
### Translation to OEL

$Sal() = Mr \leftarrow K[db][Gender() = Male]$   
 $Sal() = Ms \leftarrow K[db][Gender() = Female] \wedge K[db][MS() = Single]$   
 $Sal() = Mrs \leftarrow K[db][Gender() = Female] \wedge K[db][MS() = Married]$   
 $Sal() = Lady \leftarrow K[db][Gender() = Female] \wedge \forall m \in MS: \neg K[db][MS() = m]$   
 $Sal() = Customer \leftarrow \forall g \in G: \neg K[db][Gender() = g]$

### User knowledge

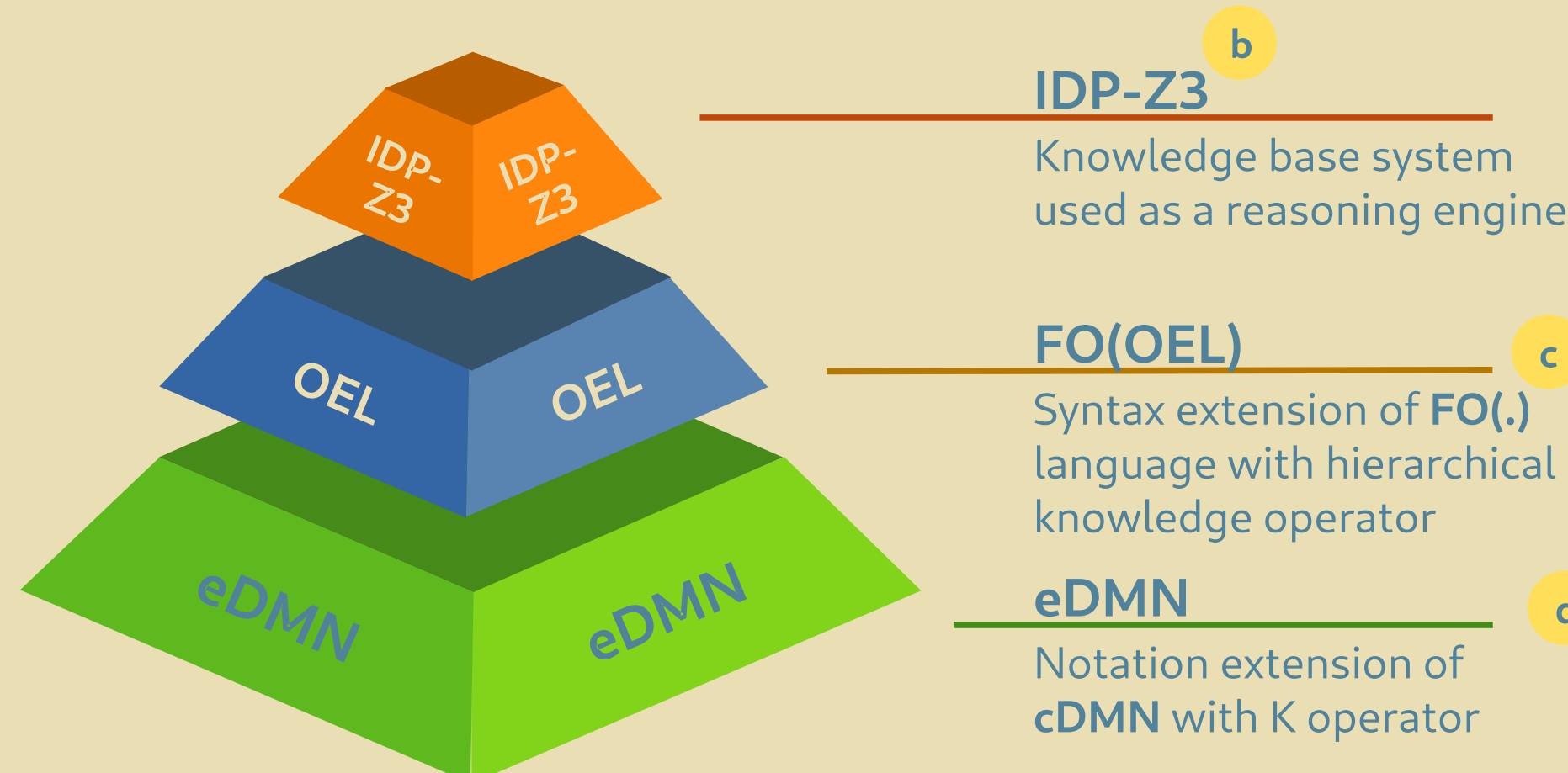


### Decision



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## Implementation<sup>a</sup>



### Developed tools:

- Ordered epistemic logic (OEL) solver
- OEL editor and Visual Studio code language support
- eDMN notation – Extension of cDMN with epistemic operators and translation to OEL

<sup>a</sup> <https://gitlab.com/krr/idp-z3-oel>

<sup>b</sup> <https://www.idp-z3.be/>

<sup>c</sup> <https://dtai.cs.kuleuven.be/pages/software/idp/FO-dot>

<sup>d</sup> <https://gitlab.com/EAVISE/cdmn/cdmn-solver>

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## Other results + Future

**Theorem 1:** Every utility function equipped with an optimization procedure and every DMN decision function can be reduced to an epistemic decision function.

**Theorem 2:** Every epistemic decision function can be modeled as an order epistemic logic theory.

**Theorem 3:** Every epistemic decision function can be modeled as an eDMN table.

### Future work:

- Generalized OEL framework for decision modeling
- Other applications for OEL solver

1

Martin Peterson. 2017. An introduction to decision theory. CUP.

2

Object Modelling Group. 2023. Decision model and notation. <https://www.omg.org/spec/DMN/>.

3

Hanne Vlaeminck, Joost Vennekens, Maurice Bruynooghe, and Marc Denecker. 2012. Ordered epistemic logic: Semantics, complexity and applications. In Thirteenth International Conference on the Principles of KRR.



Partners

